

9/9/19
Monday

HW 2

9(c) $A_n = (2 + \frac{1}{n}, n) \subseteq \mathbb{R}$

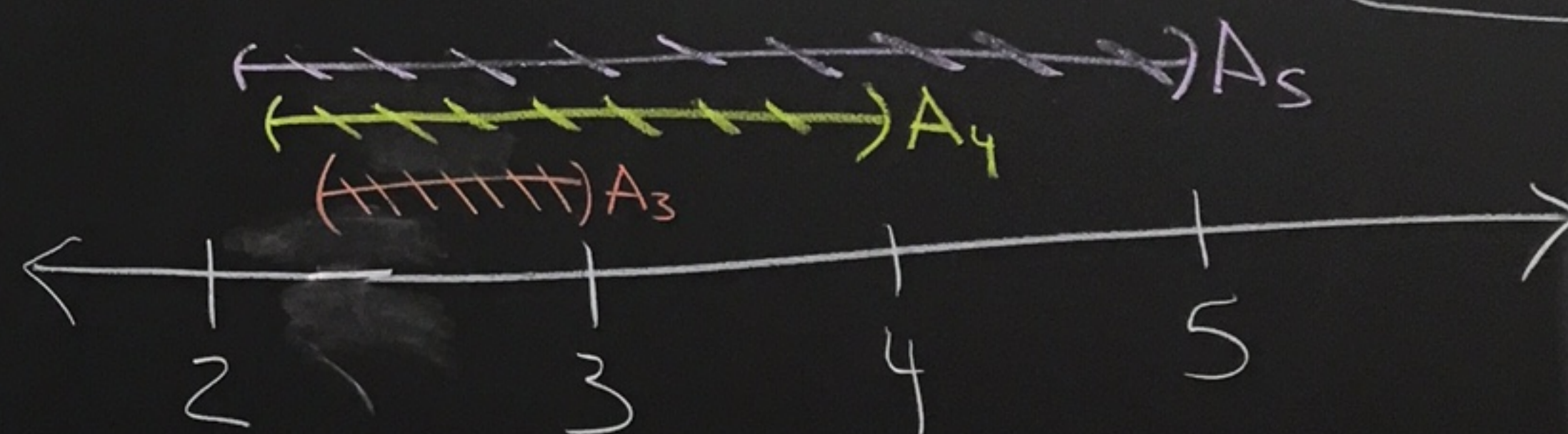
Calculate $\bigcup_{n=3}^{\infty} A_n$ and $\bigcap_{n=3}^{\infty} A_n$

$$A_3 = (2 + \frac{1}{3}, 3)$$

$$A_4 = (2 + \frac{1}{4}, 4)$$

$$A_5 = (2 + \frac{1}{5}, 5)$$

$\vdots$
 $\vdots$



$$\bigcup_{n=3}^{\infty} A_n = (2, \infty)$$

$$\bigcap_{n=3}^{\infty} A_n = A_3 = (2 + \frac{1}{3}, 3) = (\frac{7}{3}, 3)$$

From last time

Prop: Let $\mathcal{A} = \{A_\alpha \mid \alpha \in \Delta\}$

be an indexed family of sets.

Pick $\alpha_0 \in \Delta$. Then:

$$\textcircled{1} A_{\alpha_0} \subseteq \bigcup_{\alpha \in \Delta} A_\alpha$$

$$\textcircled{2} \bigcap_{\alpha \in \Delta} A_\alpha \subseteq A_{\alpha_0}$$

pf:

$\textcircled{1}$ Let $x \in A_{\alpha_0}$.

Recall that

$$\bigcup_{\alpha \in \Delta} A_\alpha = \left\{ y \mid \begin{array}{l} \text{there exists } \alpha' \in \Delta \\ \text{where } y \in A_{\alpha'} \end{array} \right\}$$

Since $x \in A_{\alpha_0}$ there does indeed exist $\alpha' \in \Delta$ with $x \in A_{\alpha'}$, i.e. $\alpha' = \alpha_0$.

So $x \in \bigcup_{\alpha \in \Delta} A_\alpha$.

Thus, $A_{\alpha_0} \subseteq \bigcup_{\alpha \in \Delta} A_\alpha$.

$\textcircled{2}$ You try.



(HW 3 TOPIC) — Equivalence relations & Well-defined operations
& Modulo n

Def: A relation $\sim$ on a set S is a subset of $S \times S$. If (x, y) is an element of $\sim$ then we write $x \sim y$ and say that x is related to y .
If (x, y) is not an element of $\sim$ then we write $x \not\sim y$ and say that x is not related to y .

Ex: $S = \{\square, \$, \phi\}$

$$\sim = \{(\square, \phi), (\square, \square), (\square, \$), (\phi, \square), (\phi, \$)\}$$

$$\square \sim \phi \text{ since } (\square, \phi) \in \sim.$$

$$\phi \sim \square \text{ since } (\phi, \square) \in \sim.$$

$$\square \sim \square \text{ since } (\square, \square) \in \sim.$$

$$\$ \not\sim \$ \text{ since } (\$, \$) \notin \sim.$$

More common way to define a relation is like the following example.

Ex: $S = \mathbb{Z}$

Define $x \sim y$ to mean $x < y$.

Then $\sim$ is a relation on $\mathbb{Z}$,

We have

$2 \sim 3$ since $2 < 3$.

$2 \not\sim -1$ since $2 \not< -1$.

→ How can we think of $\sim$ as
a subset of $\mathbb{Z} \times \mathbb{Z}$?

$$\sim = \{(x, y) \mid x < y\}$$

$$= \{(2, 3), (-100, 0), (7, 777), \dots\}$$

infinately
many
more
elements

Def: Let $\sim$ be a relation on a set S .

We say that $\sim$ is an equivalence relation on S if

① (reflexive) $x \sim x$ for all $x \in S$.

② (symmetric) If $x, y \in S$ and $x \sim y$, then $y \sim x$.

③ (transitive) If $x, y, z \in S$ and $x \sim y$ and $y \sim z$,
then $x \sim z$.

can
think
of an
equivalence
relation
as an
"equals"
sign.

Ex: Let $S = \{-15, 2, 0\}$

and $\sim = \{(0,0), (0,2), (2,0), (-15,2)\}$

(reflexive?) $0 \sim 0$ but $2 \not\sim 2$ and $(-15) \not\sim (-15)$.
So, $\sim$ is not reflexive.

(symmetric?) $\boxed{\text{NO!}}$ $(-15) \sim 2$ but $2 \not\sim (-15)$.

(transitive?) $0 \sim 0$ and $0 \sim 2 \Rightarrow 0 \sim 2$ ✓
 $0 \sim 2$ and $2 \sim 0 \Rightarrow 0 \sim 0$ ✓
 $2 \sim 0$ and $0 \sim 0 \Rightarrow 2 \sim 0$ ✓

$0 \sim 0$ and $0 \sim 0 \Rightarrow 0 \sim 0$ ✓
 $-15 \sim 2$ and $2 \sim 0$ but $-15 \not\sim 0$

So, $\sim$ is not transitive.

$\sim$ is NOT an equivalence relation

Ex: $S = \{1, 2, 3\}$

$\sim = \{(1,1), (2,2), (3,3), (1,3), (3,1)\}$

Is $\sim$ an equivalence relation?

YES

reflexive

$1 \sim 1$
 $2 \sim 2$
 $3 \sim 3$

Yes

symmetric

Yes

If $x \sim y$, then
 $y \sim x$,

Ex: $1 \sim 3$ and
 $3 \sim 1$.

transitive

- $1 \sim 1$ and $1 \sim 1$. And $1 \sim 1$.
- $1 \sim 1$ and $1 \sim 3$. And $1 \sim 3$.
- $2 \sim 2$ and $2 \sim 2$. And $2 \sim 2$.
- $3 \sim 3$ and $3 \sim 3$. And $3 \sim 3$.
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Yes

Def: Let $\sim$ be an equivalence relation on a set S . Let $x \in S$. The equivalence class of x is

$$\bar{x} = \{y \in S \mid x \sim y\}$$

$\bar{x}$ consists of all elements y that are related to x .

It doesn't matter if you define $\bar{x}$ as $\{y \in S \mid x \sim y\}$ or $\{y \in S \mid y \sim x\}$ since $\sim$ is symmetric.

Another notation for $\bar{x}$ is $[x]$.

Ex: $S = \{1, 2, 3\}$

$$\sim = \{(1,1), (2,2), (3,3), (1,3), (3,1)\}$$

$\sim$ is an equivalence relation.

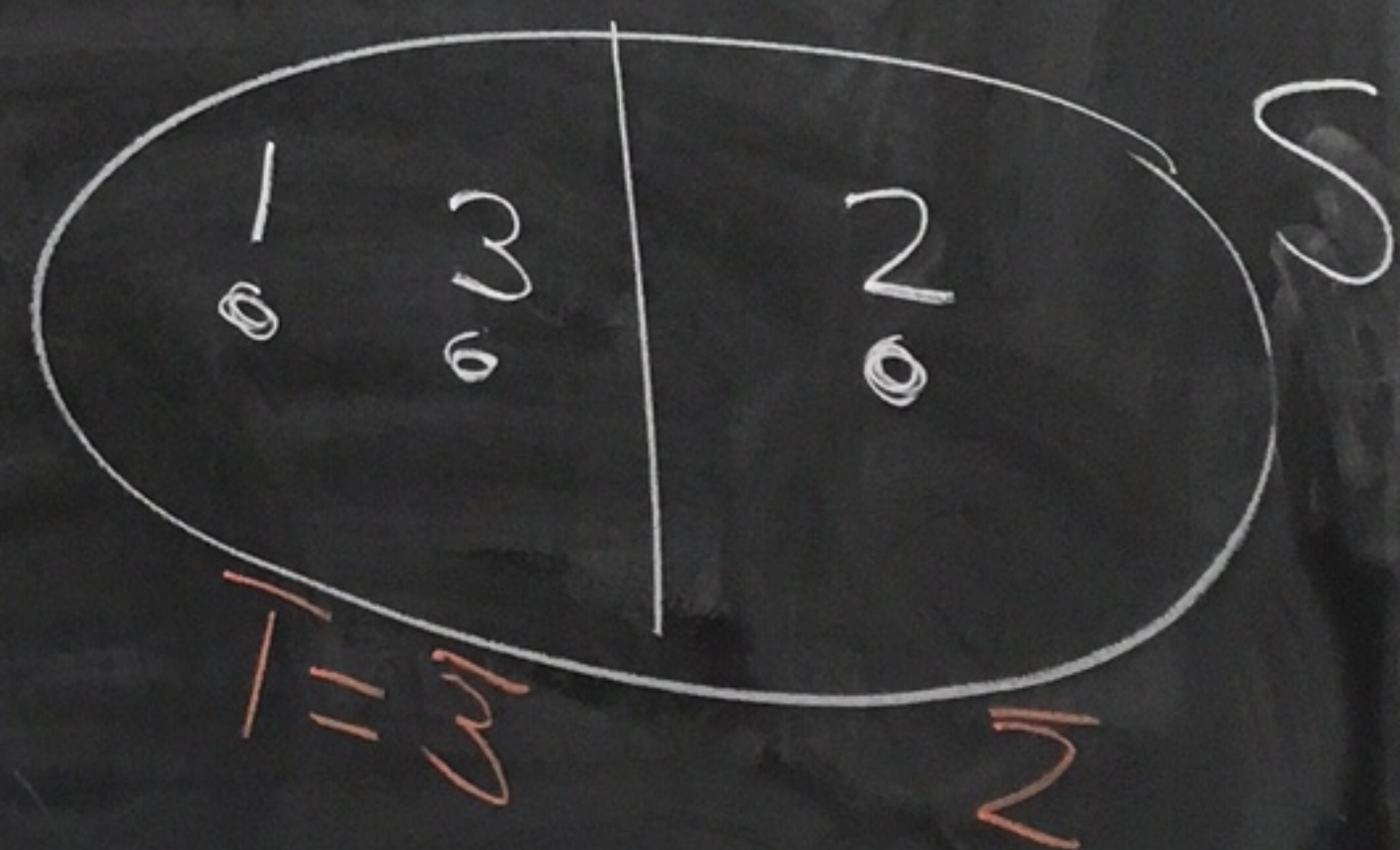
$$\overline{1} = \{y \in S \mid 1 \sim y\} = \{1, 3\}$$

$$\overline{2} = \{y \in S \mid 2 \sim y\} = \{2\}$$

$$\overline{3} = \{y \in S \mid 3 \sim y\} = \{1, 3\}$$

There are 2
equivalence classes

$$\overline{1} = \overline{3} \text{ and } \overline{2}$$



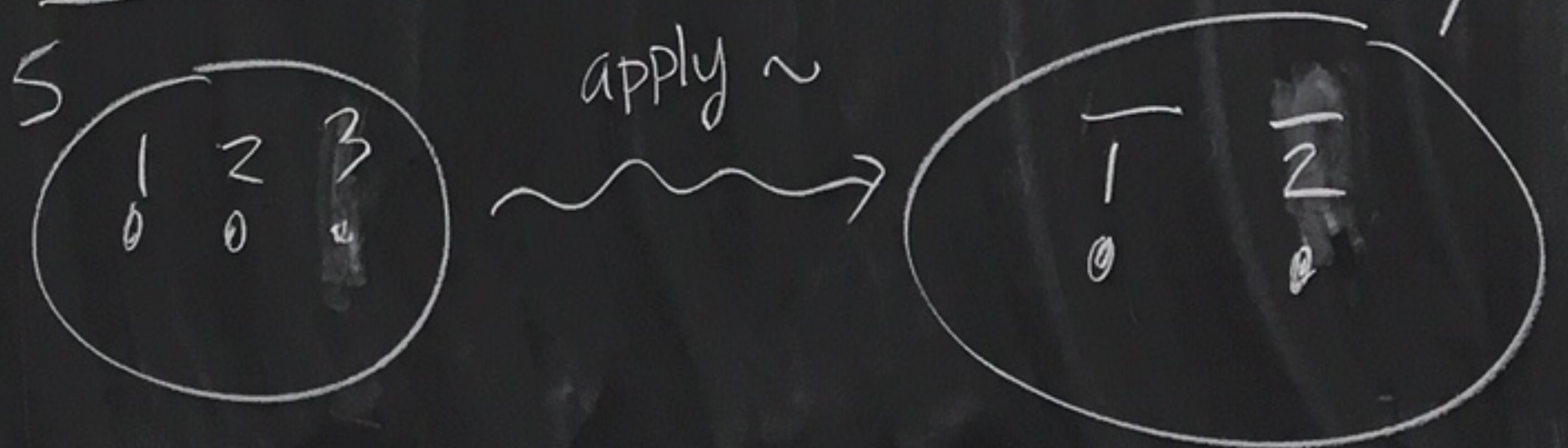
Def: Let S be a set
and $\sim$ be an equivalence
relation on S . We
denote the set of
equivalence classes
as $S/\sim$ read:
" $S \text{ mod } \sim$."

Ex: In the previous
example

$$S/\sim = \{ \overline{1}, \overline{2} \}$$

$\uparrow$
 $\overline{1} = \overline{3}$

Idea:



since $1 \sim 3$
we make
1 and 3
equal in
the set
of equivalence
classes