Math 4460 2/8/23

We want an algorithm to  
calculate gcd 
$$(a,b)$$
.  
The next theorem will be the  
basis for the Euclidean Algorithm  
Theorem: Let a and b be  
positive integers and  $0 \le a \le b$ .  
Suppose  $b = aq + r$  where  
 $r,q \in \mathbb{Z}$  and  $0 \le r \le a$ .  
Then,  
 $gcd(b,a) = gcd(a,r)$   
We replace  
this problem  
with a smaller problem

EX: Calculate gcd (138,62) 62 138 138=62(2)+14 124 Theorem says: gcd(138,62) = gcd(62,14)Kepeat the process: 1462 62 = 14(4) + 6- 56 Theorem says: 9cd(62,14) = gcd(14,6)Repeat the process: |4 = 6(2) + 2Theorem says: 9cd(14,6) = 9cd(6,2)

Repeat the process: 6 = 2(3) + 0Theorem says: g(d(6,2) = g(d(2,0))Summary: 9cd(138,62) = 9cd(62,14) = 9cd(14,6)= qcd(6,2) = qcd(2,0) = 215 Wer: gcd(138,62) = 2

proof of theorem: Let a, b EZ and O<a≤b. Use the division algorithm to Write b = aq + r with  $0 \leq r < q$ . Let d = gcd(b, a)and d' = ycd(a, r). Our goal is to show d=d. Step 1: Let's show d'≤d. Since d'=gcd(a,r) we know d'la and d'lr. So, a = d'm and r=d'n where m, nEZ.

Ergo, b = aq + r= d'mq + d'n $= q_{(md+v)}$ this is an integer because M, J, n are integers. Consequently, d'[b. Thus, d'Ib and d'Iq. So, d'is a positive common divisor of b and a. But d is the greatest pusitive

where 
$$divisor$$
 of  $b$  and  $a$ .  
Thus,  $d' \leq d$ .  
Step 2: Let's show  $d \leq d'$   
Since  $d = gcd(a,b)$  we know  
 $dla$  and  $dlb$ .  
Hence,  
 $a = ds$  and  $b = dt$   
where  $s, t \in \mathbb{Z}$ .  
It follows that  
 $r = b - aq$   
 $= dt - dsq$ 

$$= d [t - sq]$$

$$= d [t - sq]$$

$$= d [t - sq]$$

$$= d [a is an integer since t_{s,q} \in \mathbb{Z}]$$
So, d[r.
$$= d [a and d]r.$$

$$= d [a and d]r.$$
Since d'= gcd (a,r) we
$$= Know d \leq d'.$$

Therefore, since  $d \leq d$  and  $d \leq d'$ , we may conclude that d = d'.

Evelidean Algorithm (Finds gcd(b,al) Let a and b be positive integers with OKa≤b. Stepl: Divide a into b to get b = aqtcwith  $0 \leq r < \alpha$ . Step 2: If r=0, then you're done. The answer is a. If r = 0, then repeat step | but with b replaced by a and a replaced by c.

While loop  $\alpha = \#_j$ b = #;r = remainder [b, a]; While (r=+0)  $b = \alpha_j$  $\alpha = (\hat{})$ r=remainder [b,a]; Print [ n ];

Recursion method  $qcd(b, \alpha) := [$ r=remainder b,a); TF | r=0return (a); else return [gcd(a,r)]; 

Ex: Find gcd (578,153)

 $578 = 3 \cdot 153 + 119$   $153 = 1 \cdot 119 + 34$   $119 = 3 \cdot 34 + 17$   $34 = 2 \cdot 17 + 07$ Answer 9cd(578, 153) = 17

g(d(578,153))= g(d(153,119))= g(d(119,34))= g(d(34,17))= g(d(17,0))= 17



So, I<a<n since all. Since aln we know n = ab where  $b \in \mathbb{Z}$ . Since a &n me positive, so is b. We have  $b = \frac{\pi}{\alpha}$ . Then,  $| < \frac{n}{a} = b$ because a < hAnd,  $b = \frac{n}{a} < n$  because 1< So, | < b < n, | < a > b < n $S_{o}$  | < b < n. Sp, N = ab where | < a < n, | < b < n.  $\overline{}$